

Lecture 2: DS 402: Big Data Concepts

HDFS, Inside Map-Reduce, Introduction to Machine Learning

Outline

- Hadoop Distributed File System
- Recap SQL, Data Structure and Programming
- HashMaps
- Inside Map-Reduce-Twitter Word Count
- Introduction to Machine Learning
- Naive Bayes

Lab/Assignment
Spam-Ham Problem

No arbitrary file modifications!

Files are modified by **appending new data** rather than overwriting existing data. Random writes within a file are practically non-existent. Once written, the files are only read, and often only sequentially.

Data Flow

** Single Master

- However, we must minimize its involvement in reads and writes so that it does not become a bottleneck.
- Clients never read and write file data through the master.
- Instead, a client asks the master which chunk-servers it should contact. It caches

The Google File System

Sergey Brin, Michael G. Christl, Shun-Tae Lee, The Google file system, Proceedings of the International ACM Symposium on Operating Systems Principles, October 19-22, 2003, Boston, MA, USA



Distributed File System

** Fault-tolerance at scale

- application bugs
- operating system bugs
- human errors
- failures of disks, memory
- connectors
- networking
- power supplies

Therefore, constant monitoring, error detection, built-in recovery, and automatic recovery must be integral to the system.

The Design of GFS/HDFS

Assumptions

In designing a file system for our needs, we have been guided by assumptions that offer both challenges and opportunities.

- The system is built from many **inexpensive commodity** components that often fail. It must constantly monitor itself and detect, tolerate, and recover promptly from component failures on a routine basis.
- The system stores a modest number of large files. We expect a few million files, each typically 100 MB or larger in size. Multi-GB files are the common case and should be managed efficiently. Small files must be supported, but we need not optimize for them. The workloads primarily consist of few kinds of reads: large streaming reads and small random reads.
- The workloads also have many large, sequential writes that append data to files.

Block / Chunk Size

A disk has a block size, which is the minimum amount of data that it can read or write.
512 bytes
64 MB

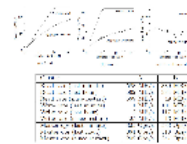
Minimize the cost of seeks.
Thus the drive to optimize a large file reads multiple blocks operates at the disk transfer rate.

A disk's random access rate is the rate at which it can read or write data. This is much slower than the disk transfer rate.

The disk is actually 64 MB, although many GFS implementations use 128 MB blocks. The reason will probably be to accommodate for transfer speeds that will have gone up to 4 GB/s.

File system metadata is usually stored on a separate disk or partition. Some implementations store metadata in the blocks, just prior to the first data block.

Read and Writes Rates

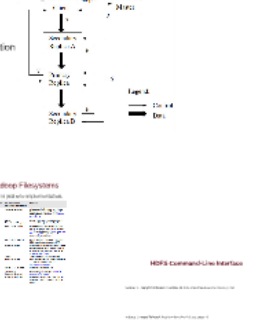


The Word-Count Example



Data Redundancy

- Mechanism
- Re-replication
 - Re-balancing
 - Garbage Collection



Recap-SQL, DS and Basic Programming

City	Area	Population
New York	1000	10000000
Los Angeles	1000	10000000
Chicago	1000	10000000
San Francisco	1000	10000000
London	1000	10000000
Paris	1000	10000000
Madrid	1000	10000000
Rome	1000	10000000
Beijing	1000	10000000
Mumbai	1000	10000000
Delhi	1000	10000000
Shanghai	1000	10000000
Manila	1000	10000000
Seoul	1000	10000000
Osaka	1000	10000000
Tokyo	1000	10000000

SQL: SELECT SUM(Pop) FROM Table GROUP BY City

** HBase and Hive

Recap Data Structures and Programming

- HashMaps
- Java
- Sorting
- Searching

Why Machine Learning?

- Big Data - The Problem
- Volume: Terabytes
 - Variety: Structured, Unstructured
 - Velocity: Real-time
 - Veracity: Inconsistent
 - Variability: Changing



Why Machine Learning? (continued...)

- Spam or Not?
- Yes
 - No
- Web Page Categorization
- News
 - Sports
 - Misc
 - Blog
 - Product
 - ...

** Impossible to write a concise rule-set!

Naive Bayes Algorithm

Naive Bayes

- Supervised
- Naive = Simple
- Simple representation of document
- Bag of words

Bag-of-Words Representation

$$Y = \begin{pmatrix} \text{word}_1 & \text{word}_2 & \dots & \text{word}_n \\ \text{count}_1 & \text{count}_2 & \dots & \text{count}_n \end{pmatrix} = C$$

NB-Mathematical Foundation

$$P(C) = \prod_{i=1}^n P(\text{word}_i)$$

$$P(C) = \prod_{i=1}^n \frac{c_i}{N}$$

NB-Working Example

Word	Count	Probability
the	10	0.1
is	5	0.05
and	3	0.03
of	2	0.02
to	1	0.01
in	1	0.01
on	1	0.01
with	1	0.01
at	1	0.01
by	1	0.01
from	1	0.01
as	1	0.01
for	1	0.01
with	1	0.01
the	1	0.01
of	1	0.01
and	1	0.01
is	1	0.01
to	1	0.01
in	1	0.01
on	1	0.01
with	1	0.01
at	1	0.01
by	1	0.01
from	1	0.01
as	1	0.01
for	1	0.01

Assignment

Assignment: Implement Naive Bayes Classifier

Training

- Task Set 1
- Learn prior probabilities: spam and ham
 - Stop Words
 - Laplace Smoothing
 - Spam or Not? Tagging
 - Spamming

Testing

- Given an unknown text, judge automatically whether it is a ham or spam
 - Run on Set 2
- Measuring System's Accuracy
- Accuracy (ACC) is an overall score. For any given classification task, the accuracy is the ratio of the number of correct classifications to the total number of classifications.
- Accuracy (ACC) = (Number of correct classifications) / (Total number of classifications)

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Block / Chunk Size

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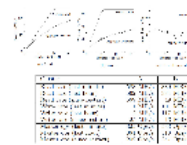
Minimize the cost of seeks.
Thus the drive to optimize a large file reads multiple blocks operates at the disk transfer rate.

A disk controller reads 8K of the disk block at a time, and the transfer rate is 100 MB/s. That is, it can read the 256 of the data that it can read in one block at a time.

The disk is actually 64 KB, although many GFS implementations use 128 KB blocks. The reason will probably be to accommodate the transfer speed (up with some number of disk heads).

File system metadata is usually stored on a separate disk. If you have two disk heads, there will be two heads in the disk, and you will get better performance.

Read and Writes Rates

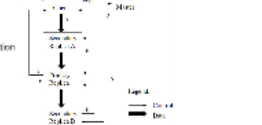


The Word-Count Example



Data Redundancy

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Hadoop FileSystems
HDFS is a distributed file system designed to run on commodity hardware. It is designed to store very large files across many machines, and to support high throughput retrieval of file data.

HDFS Command-Line Interface

Recap-SQL, DS and Basic Programming

City	Population	Area
New York	18,000,000	30,000
Los Angeles	12,000,000	20,000
Chicago	10,000,000	15,000
San Francisco	8,000,000	10,000
London	10,000,000	15,000
Paris	10,000,000	15,000
Madrid	6,000,000	10,000
Rome	6,000,000	10,000
Beijing	18,000,000	30,000
Mumbai	18,000,000	30,000
Delhi	18,000,000	30,000
London	10,000,000	15,000
Paris	10,000,000	15,000
Madrid	6,000,000	10,000
Rome	6,000,000	10,000
Beijing	18,000,000	30,000
Mumbai	18,000,000	30,000
Delhi	18,000,000	30,000

SQL: SELECT SUM(Population) FROM Table GROUP BY City

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 - Variety: Structured, Unstructured
 - Velocity: Real-time data
 - Value: Machine Learning



Why Machine Learning? (continued...)

- Spam or Not?
- News
 - Sports
 - Wine
 - Blog
 - Product

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What is Machine Learning?

Machine Learning is the study of algorithms that learn from data and use the learned information to solve a problem.

ML Types

Naive Bayes Algorithm

Naive Bayes

- Supervised
- Naive = Simple
- Simple representation of document
- Bag of words

Bag-of-Words Representation

$$Y = \begin{pmatrix} \text{word}_1 & \text{word}_2 & \dots & \text{word}_n \\ \text{count}_1 & \text{count}_2 & \dots & \text{count}_n \end{pmatrix} = C$$

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$$P(C) = \prod_{i=1}^n P(\text{word}_i)$$

$$P(C) = \prod_{i=1}^n P(\text{word}_i | \text{word}_1, \dots, \text{word}_{i-1})$$

NB-Working Example

Word	Count
spam	1
ham	1
word1	1
word2	1
word3	1
word4	1
word5	1
word6	1
word7	1
word8	1
word9	1
word10	1
word11	1
word12	1
word13	1
word14	1
word15	1
word16	1
word17	1
word18	1
word19	1
word20	1

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 - HashMap
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The Google File System

Sanjay Ghemawat , Howard Gobioff , Shun-Tak Leung, The Google file system, Proceedings of the nineteenth ACM symposium on Operating systems principles, October 19-22, 2003, Bolton Landing, NY, USA



Sanjay Ghemawat



Howard Gobioff



Shun-Tak Leung

The Google File System

Sanjay Ghemawat, Howard Gobioff, and Shun-Tak Leung
Google*

ABSTRACT

We have designed and implemented the Google File System, a scalable distributed file system for large distributed data-intensive applications. It provides **fault tolerance** while running on inexpensive **commodity hardware**, and it delivers high aggregate performance to a large number of clients.

While sharing many of the same goals as previous distributed file systems, our design has been driven by observations of our application workloads and technological environment, both current and anticipated, that reflect a marked departure from some earlier file system assumptions. This

1. INTRODUCTION

We have designed and implemented the Google File System (GFS) to meet the rapidly growing demands of Google's data processing needs. GFS shares many of the same goals as previous distributed file systems such as performance, scalability, reliability, and availability. However, its design has been driven by key observations of our application workloads and technological environment, both current and anticipated, that reflect a marked departure from some earlier file system design assumptions. We have reexamined traditional choices and explored radically different points in the

Distributed File System

**** *Fault-tolerance at scale***

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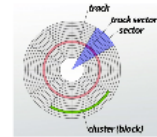
Files are mutated by *appending new data rather than overwriting existing data*. Random writes within a file are practically non-existent. Once written, the files are only read, and often only *sequentially*.

Block / Chunk Size

A disk has a block size, which is the minimum amount of data that it can read or write.

512 bytes

64 MB



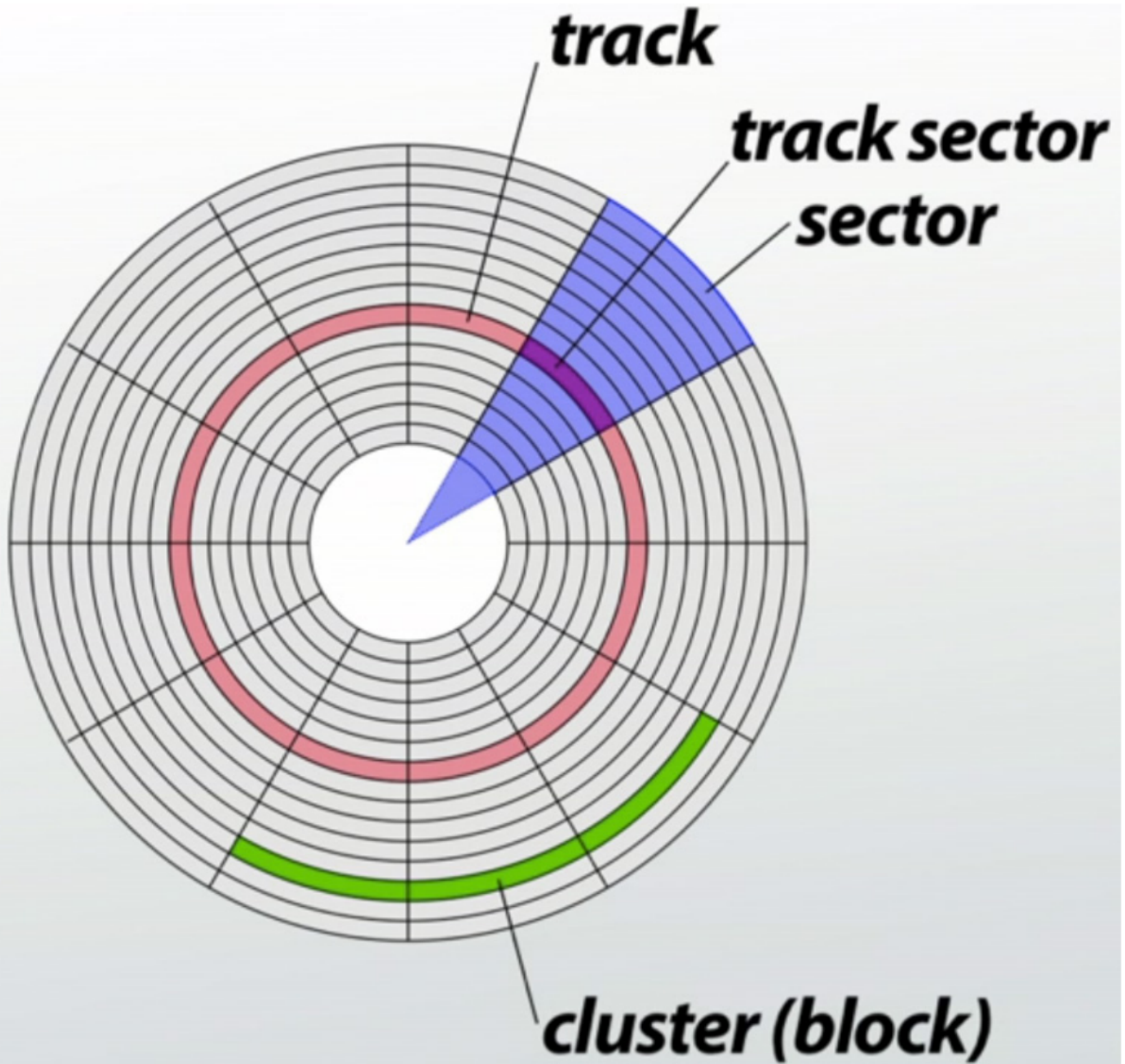
minimize the cost of seeks.

Thus the time to transfer a large file made of multiple blocks operates at the disk transfer rate.

A quick calculation shows that if the seek time is around 10 ms, and the transfer rate is 100 MB/s, then to make the seek time 1% of the transfer time, we need to make the block size around 100 MB.

The default is actually 64 MB, although many HDFS installations use 128 MB blocks. This figure will continue to be revised upward as transfer speeds grow with new generations of disk drives.

** Map tasks in MapReduce normally operate on one block at a time, so if you have too few tasks (fewer than nodes in the cluster), your jobs will run slower than they could otherwise.

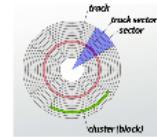


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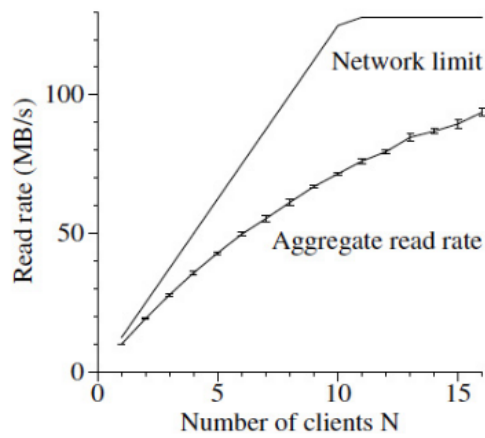
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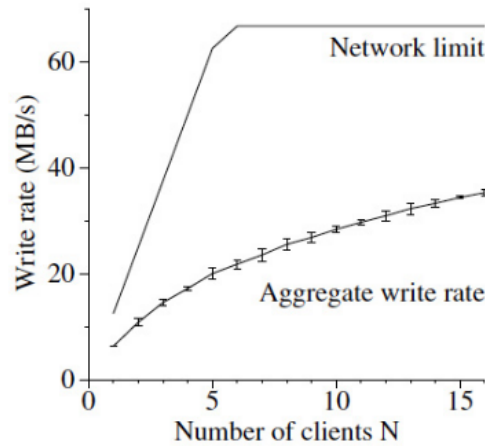
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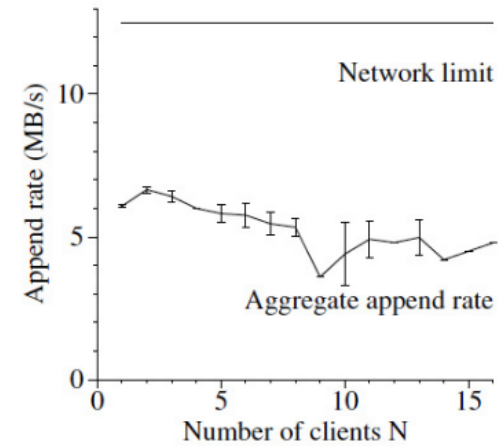
Read and Writes Rates



(a) Reads



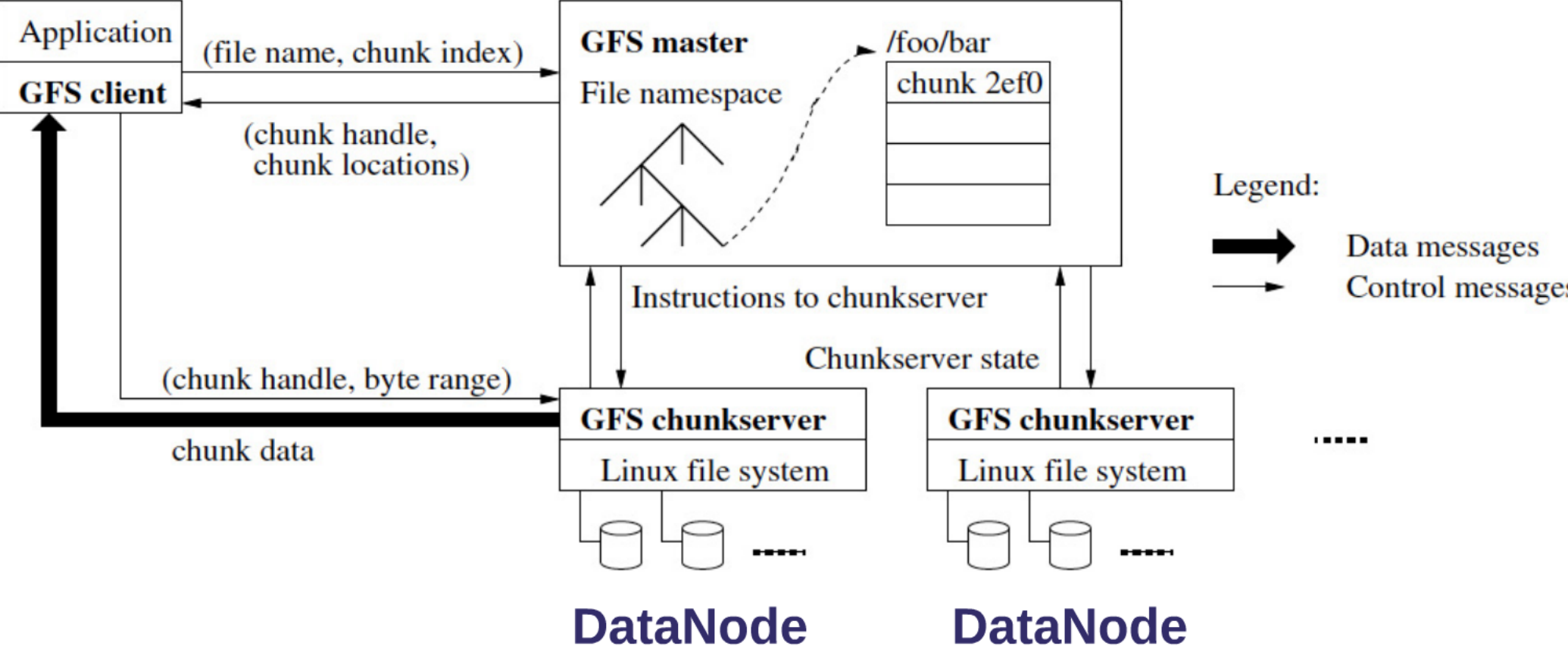
(b) Writes



(c) Record appends

Cluster	A	B
Read rate (last minute)	583 MB/s	380 MB/s
Read rate (last hour)	562 MB/s	384 MB/s
Read rate (since restart)	589 MB/s	49 MB/s
Write rate (last minute)	1 MB/s	101 MB/s
Write rate (last hour)	2 MB/s	117 MB/s
Write rate (since restart)	25 MB/s	13 MB/s
Master ops (last minute)	325 Ops/s	533 Ops/s
Master ops (last hour)	381 Ops/s	518 Ops/s
Master ops (since restart)	202 Ops/s	347 Ops/s

NameNode





Shun-Tak Leung

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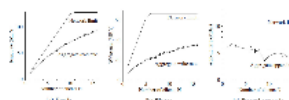
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Therefore, constant monitoring, error detection, fault tolerance, and automatic recovery must be integral to the system.

Read and Writes Rates



Cluster	A	B
Read rate (MB/s)	185.5 MB/s	303.5 MB/s
Write rate (MB/s)	292.5 MB/s	304.5 MB/s
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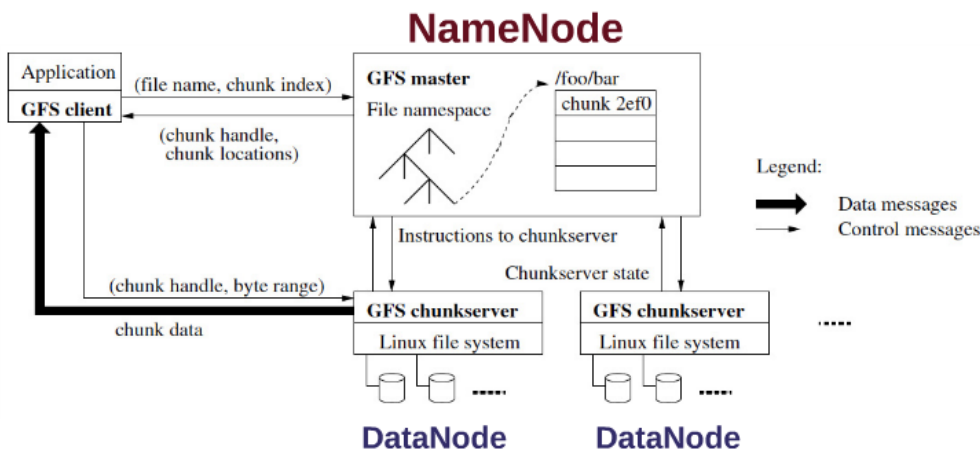
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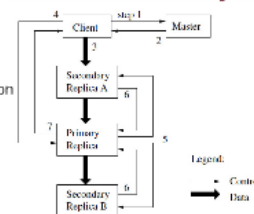
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Data Redundancy

Mechanism

- Re-replication
- Re-balancing
- Garbage Collection



Hadoop Filesystems
HDFS is just one implementation.

HDFS Command-Line Interface

Recap-SQL, DS and Basic Programming

City	Sales
Beijing	1180000
Chengde	1150000
Chongqing	1220000
Guangzhou	1400000
Nanjing	1700000
Shanghai	1700000

SQL: SELECT SUM(Sales) from Table GROUP BY City

** HBase and Hive

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Introduction to Machine Learning

Why Machine Learning?

- Big Data - The Problem
 - Robust Technology
 - Hadoop
 - Intelligent Analytic
 - Unstructured data
 - Machine Learning



Why Machine Learning? (continue...)

- Spam or Ham
- Web Page Categorization
 - News
 - Sports
 - Wiki
 - Blog
 - Product
 - ...

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What is Machine Learning?

Definition 1
Arthur Samuel (1959) Machine Learning: "Part of study that gives computers the ability to learn without being explicitly programmed."

Definition 2
Tom Mitchell (1996) Artificial Learning Problem: "A computer program is said to learn from experience E, with respect to some task T, and some performance measure P, if its performance on T, as measured by P, improves with experience E."

ML Types

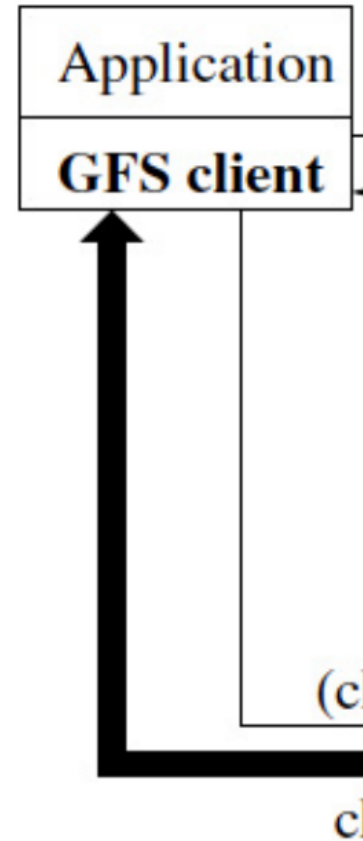
- Supervised Learning
- Unsupervised Learning
- Reinforcement Learning

Naive Bayes Algorithm

Data Flow

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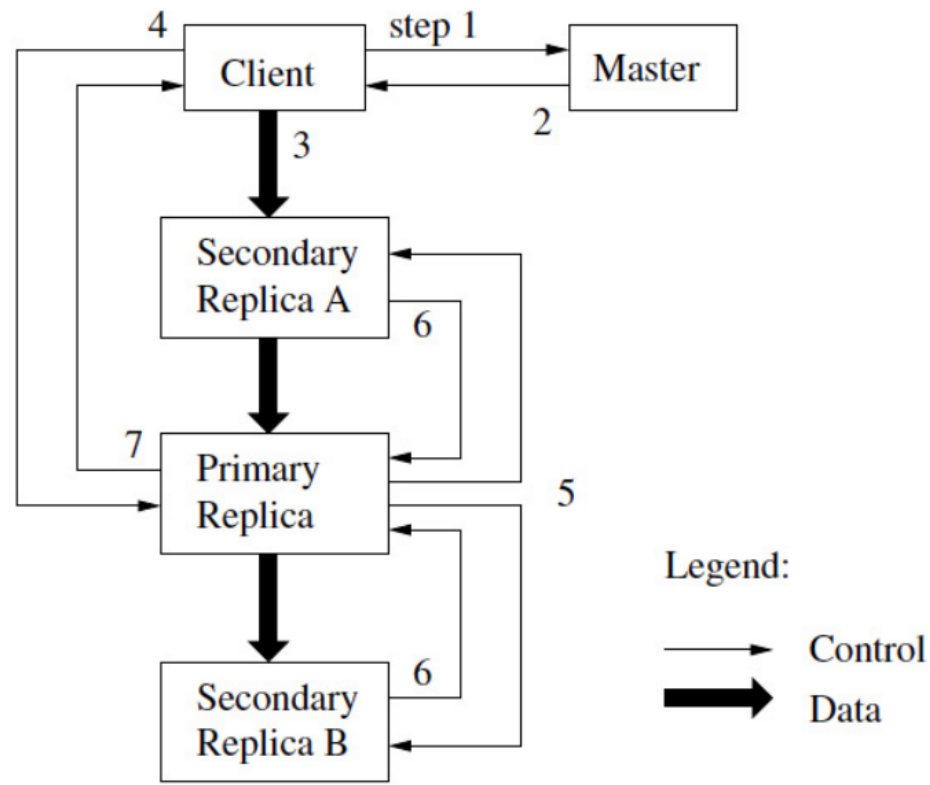
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Hadoop Filesystems

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Filesystem	URI scheme	Java implementation (all under org.apache.hadoop)	Description
Local	<i>file</i>	<code>fs.LocalFileSystem</code>	A filesystem for a locally connected disk with client-side checksums. Use <code>RawLocalFileSystem</code> for a local filesystem with no checksums. See " LocalFileSystem " on page 76.
HDFS	<i>hdfs</i>	<code>hdfs.DistributedFileSystem</code>	Hadoop's distributed filesystem. HDFS is designed to work efficiently in conjunction with MapReduce.
HFTP	<i>hftp</i>	<code>hdfs.HftpFileSystem</code>	A filesystem providing read-only access to HDFS over HTTP. (Despite its name, HFTP has no connection with FTP.) Often used with <i>distcp</i> (see " Parallel Copying with distcp " on page 70) to copy data between HDFS clusters running different versions.
HSFTP	<i>hsftp</i>	<code>hdfs.HsftpFileSystem</code>	A filesystem providing read-only access to HDFS over HTTPS. (Again, this has no connection with FTP.)
HAR	<i>har</i>	<code>fs.HarFileSystem</code>	A filesystem layered on another filesystem for archiving files. Hadoop Archives are typically used for archiving files in HDFS to reduce the namenode's memory usage. See " Hadoop Archives " on page 71.
KFS (Cloud-Store)	<i>kfs</i>	<code>fs.kfs.KosmosFileSystem</code>	CloudStore (formerly Kosmos filesystem) is a distributed filesystem like HDFS or Google's GFS, written in C++. Find more information about it at http://kosmosfs.sourceforge.net/ .
FTP	<i>ftp</i>	<code>fs.ftp.FTPFileSystem</code>	A filesystem backed by an FTP server.
S3 (native)	<i>s3n</i>	<code>fs.s3native.NativeS3FileSystem</code>	A filesystem backed by Amazon S3. See http://wiki.apache.org/hadoop/AmazonS3 .
S3 (block-based)	<i>s3</i>	<code>fs.s3.S3FileSystem</code>	A filesystem backed by Amazon S3, which stores files in blocks (much like HDFS) to overcome S3's 5 GB file size limit.

HDFS Command-Line Interface

```
hadoop fs -copyFromLocal input/abc.txt hdfs://localhost/user/amitava/xyz.txt
```

```
hadoop fs -copyToLocal /user/amitava/xyz.txt abc.copy.txt
```


Recap-SQL, DS and Basic Programming

City	Sales
Delhi	₹50000
Kolkata	₹40000
Gurgaon	₹30000
Bangalore	₹80000
Mumbai	₹70000
Delhi	₹55000
Bangalore	₹88000
...	...
...	...

City	Sales
Bangalore	₹168000
Delhi	₹105000
Gurgaon	₹30000
Kolkata	₹40000
Mumbai	₹70000
...	...
...	...

SQL: SELECT SUM(Sales) from Table GROUP BY City

**** HBase and Hive**

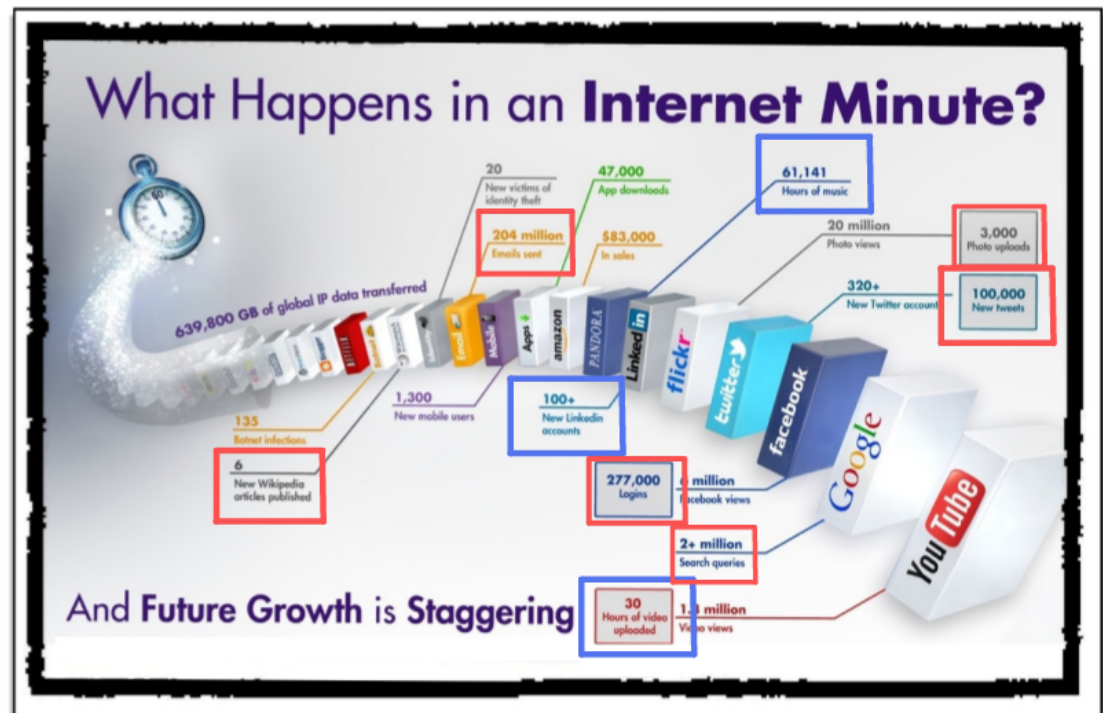
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- Searching

Why Machine Learning?

Big Data - The Problem

- Robust Technology
 - Hadoop
- Intelligent Analytic
 - Unstructured data
 - Machine Learning



What Happens in an Internet Minute?



And Future Growth is Staggering

Why Machine Learning? (continue..)

Spam or Ham



Web Page Categorization

- News
- Sports
- Wiki
- Blog
- Product
- ...

**** Impossible to write a concise rule-set!!**

What is Machine Learning?

Definition 1

Arthur Samuel (1959). Machine Learning: Field of study that gives computers the ability to learn without being explicitly programmed.

Definition 2

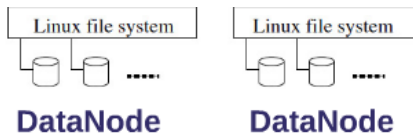
Tom Mitchell (1998) Well-posed Learning Problem: A computer program is said to learn from experience E with respect to some task T and some performance measure P , if its performance on T , as measured by P , improves with experience E .

ML Types

- Supervised learning
- Unsupervised learning
- Online Learning
- Others:
 - Reinforcement learning

master which chunk-servers it should contact. It caches

chunk data



Recap-SQL, DS and Basic Programming

City	Sales
Bangalore	7166000
Chennai	2102000
Gurgaon	1311000
Kolkata	8631000
Mumbai	7311000
...	...

SQL: SELECT SUM(Sales) from Table GROUP BY City
 ** HBase and Hive

Recap Data Structures and Programming

- HashMaps
- Java
- Sorting
- Searching

Why Machine Learning?

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Introduction to Machine Learning

Why Machine Learning? (continue..)

Spam or Ham



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Web Page Categorization

- News
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What is Machine Learning?

Definition 1
 John Elomri (2006). Machine Learning: Principles and Applications. Morgan Kaufmann Publishers, Inc. ISBN 0-12-812181-7.

Definition 2
 Tom Mitchell (2004). Machine Learning: Probabilistic and Statistical Approaches. Morgan Kaufmann Publishers, Inc. ISBN 0-12-812181-7.

ML Types

- Supervised Learning
- Unsupervised Learning
- Reinforcement Learning
- Deep Learning

Naive Bayes Algorithm

Naive Bayes

- Supervised
- "Naive" = Simple
- simple representation of document
- Bag of words

Bag-of-Words Representation

$$Y(x) = \begin{pmatrix} \text{great} & 2 \\ \text{love} & 2 \\ \text{excited} & 1 \\ \text{length} & 1 \\ \text{happy} & 1 \\ \dots & \dots \end{pmatrix} = C$$

NB-Mathematical Foundation

For a document d and c class:
 $c_{d,c} = \sum_{i=1}^n P(x_i | c)$
 $= \arg \max_c P(x_i | c)$
 $= \arg \max_c \prod_{i=1}^n P(x_i | c)$
 $= \arg \max_c \prod_{i=1}^n P(x_i | c) P(c)$
 Conditional independence: assume the features (words) are independent given the class c .

NB-Working Example

Class	great	love	excited	length	happy
Spam	0.000000	0.000000	0.000000	0.000000	0.000000
Ham	0.000000	0.000000	0.000000	0.000000	0.000000

Assignment

<http://www.cse.cmu.edu/~jerryzhu/data/naive-spam/>

- Stop Words
- Part-of-Speech tagging
- Stemming

Training

Take Set 1

- Learn prior polarities: **spam and ham**
- Stop Words.
- Part-of-Speech tagging
- Stemming

Testing

- Given an unknown mail: judge automatically whether it is a ham or spam
- Run on Set 2

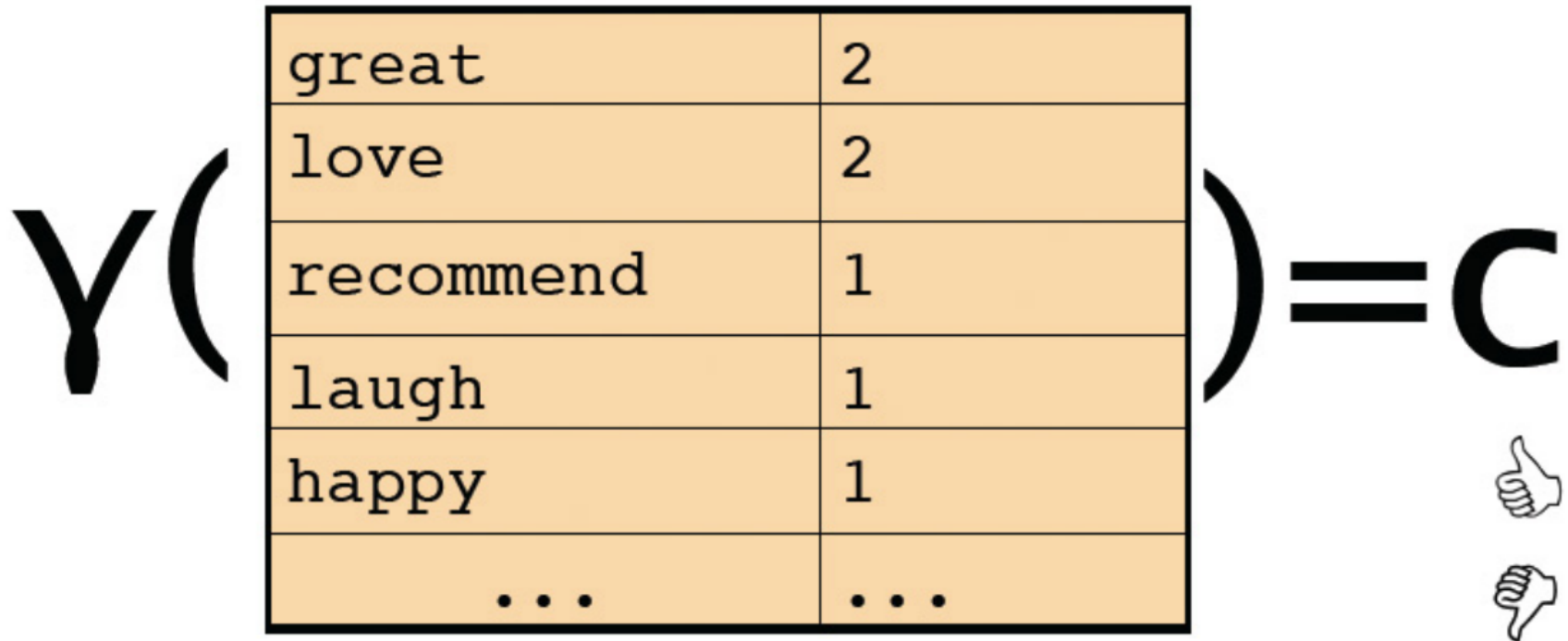
Measuring System's Accuracy

Accuracy = correctly identified/total instances
 Precision (P): % of selected items that are correct
 Recall (R): % correct items that are selected
 F-Measure (F): $2PR / (P+R)$

Naive Bayes

- Supervised
- “Naïve” = Simple
- simple representation of document
 - Bag of words

Bag-of-Words Representation



NB-Mathematical Foundation

For a document d and a class c

$$c_{MAP} = \operatorname{argmax}_{c \in \mathcal{C}} P(c | d)$$

$$= \operatorname{argmax}_{c \in \mathcal{C}} \frac{P(d | c)P(c)}{P(d)}$$

$$= \operatorname{argmax}_{c \in \mathcal{C}} P(d | c)P(c)$$

$$= \operatorname{argmax}_{c \in \mathcal{C}} P(x_1, x_2, \dots, x_n | c)P(c)$$

Prior Polarity



Conditional Independence: Assume the feature probabilities $P(x_i | c)$ are independent given the class c .

$$P(x_1, \dots, x_n | c) = P(x_1 | c) \cdot P(x_2 | c) \cdot P(x_3 | c) \cdot \dots \cdot P(x_n | c)$$

NB-Working Example

$$P(c) = \frac{N_c}{N}$$

$$P(w|c) = \frac{\text{count}(w,c)+1}{\text{count}(c)+|V|}$$

Priors:

$$P(c) = \frac{3}{4}$$

$$P(j) = \frac{1}{4}$$

	Doc	Words	Class
Training	1	Chinese Beijing Chinese	c
	2	Chinese Chinese Shanghai	c
	3	Chinese Macao	c
	4	Tokyo Japan Chinese	j
Test	5	Chinese Chinese Chinese Tokyo Japan	?

Conditional Probabilities:

$$P(\text{Chinese}|c) = \frac{(5+1)}{(8+6)} = \frac{6}{14} = \frac{3}{7}$$

$$P(\text{Tokyo}|c) = \frac{(0+1)}{(8+6)} = \frac{1}{14}$$

$$P(\text{Japan}|c) = \frac{(0+1)}{(8+6)} = \frac{1}{14}$$

$$P(\text{Chinese}|j) = \frac{(1+1)}{(3+6)} = \frac{2}{9}$$

$$P(\text{Tokyo}|j) = \frac{(1+1)}{(3+6)} = \frac{2}{9}$$

$$P(\text{Japan}|j) = \frac{(1+1)}{(3+6)} = \frac{2}{9}$$

Choosing a class:

$$P(c|d5) \propto \frac{3}{4} * \left(\frac{3}{7}\right)^3 * \frac{1}{14} * \frac{1}{14} \approx 0.0003$$

$$P(j|d5) \propto \frac{1}{4} * \left(\frac{2}{9}\right)^3 * \frac{2}{9} * \frac{2}{9} \approx 0.0001$$

Assignment

<http://www.aueb.gr/users/ion/data/enron-spam/>

- [readme.txt](#)
- Enron-Spam in pre-processed form:
 - [Enron1](#) → Spam
 - [Enron2](#) → Ham
 - [Enron3](#)
 - [Enron4](#)
 - [Enron5](#)
 - [Enron6](#)
- Enron-Spam in raw form:
 - ham messages:
 - [beck-s](#)
 - [farmer-d](#)
 - [kaminski-v](#)
 - [kitchen-l](#)
 - [lokay-m](#)
 - [williams-w3](#)
 - spam messages:
 - [BG](#)
 - [GP](#)
 - [SH](#)

Stop Words - Unmeaningful Words

a	been	doesn't	he	i'm	myself	own
about	before	doing	he'd	i've	no	same
above	being	don't	he'll	if	nor	shan't
after	below	down	he's	in	not	she
again	between	during	her	into	of	she'd
against	both	each	here	is	off	she'll
all	but	few	here's	isn't	on	she's
am	by	for	hers	it	once	should
an	can't	from	herself	it's	only	shouldn't
and	cannot	further	him	its	or	so
any	could	had	himself	itself	other	some
are	couldn't	hadn't	his	let's	ought	such
aren't	did	has	how	me	our	than
as	didn't	hasn't	how's	more	ours	that
at	do	have	i	most	ourselves	that's
be	does	haven't	i'd	mustn't	out	the
because	doesn't	having	i'll	my	over	their

POS Tagging

I hope this'll show the server working.

I/PRP hope/VBP this/DT 'll/MD show/VB the/DT server/NN working/VBG ./.

Content Words: Noun, Verb, Adjective, and Adverbs

Tag	Description
CC	Coordinating conjunction
CD	Cardinal number
DT	Determiner
EX	Existential <i>there</i>
FW	Foreign word
IN	Preposition or subordinating conjunction
JJ	Adjective
JJR	Adjective, comparative
JJS	Adjective, superlative
LS	List item marker
MD	Modal
NN	Noun, singular or mass
NNS	Noun, plural
NNP	Proper noun, singular
NNPS	Proper noun, plural
PDT	Predeterminer
POS	Possessive ending
PRP	Personal pronoun

Tag	Description
PRP\$	Possessive pronoun
RB	Adverb
RBR	Adverb, comparative
RBS	Adverb, superlative
RP	Particle
SYM	Symbol
TO	<i>to</i>
UH	Interjection
VB	Verb, base form
VBD	Verb, past tense
VBG	Verb, gerund or present participle
VBN	Verb, past participle
VBP	Verb, non-3rd person singular present
VBZ	Verb, 3rd person singular present
WDT	Wh-determiner
WP	Wh-pronoun
WP\$	Possessive wh-pronoun
WRB	Wh-adverb

Surface Forms

car, cars, car's, cars'

Stems

Car

**** Stems are not root!!**

goes -> goe

happily -> happili

Testing

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- Run on Set 2

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Precision (P): % of selected items that are correct

Recall (R): % correct items that are selected

F-Measure (F): $2PR / (P+R)$

Today in the Lab: Twitter Word Counts

- Simple Word Count
- Stop Words:
 - <http://www.ranks.nl/stopwords>
- Part-of-Speech tagging
 - <http://www.ark.cs.cmu.edu/TweetNLP/>
- Stemming
 - http://ir.dcs.gla.ac.uk/resources/linguistic_utils/porter.java